

Templates and generic programming

Improving on last week's assignment

- ◆ How did you calculate the machine precision?
 - ◆ Did you just have a main() function
 - ◆ Did you have three functions with different names?

```
epsilon_float()
epsilon_double()
epsilon_long_double()
```
 - ◆ Did you have three functions with the same name?

```
epsilon(float x)
epsilon(double x)
epsilon(long double x)
```
 - ◆ Or did you have just one function that could be used for any type?

```
epsilon()
```

Generic algorithms versus concrete implementations

- ◆ Algorithms are usually very generic:
for min() all that is required is an order relation "<"

$$\min(x,y) = \begin{cases} x & \text{if } x < y \\ y & \text{otherwise} \end{cases}$$

- ◆ Most programming languages require concrete types for the function definition

- ◆ C:

```
int min_int(int a, int b) { return a<b ? a : b;}
float min_float (float a, float b) { return a<b ? a : b;}
double min_double (double a, double b) { return a<b ? a : b;}
...
```

- ◆ Fortran:

```
MIN(), AMIN(), DMIN(), ...
```

Function overloading in C++

- ◆ solves one problem immediately: we can use the same name

```
int min(int a, int b) { return a<b ? a : b;}
float min (float a, float b) { return a<b ? a : b;}
double min (double a, double b) { return a<b ? a : b;}

```

- ◆ Compiler chooses which one to use

```
min(1, 3); // calls min(int, int)
min(1., 3.); // calls min(double, double)
```

- ◆ However be careful:

```
min(1, 3.1415927); // Problem! which one?
min(1., 3.1415927); // OK
min(1, int(3.1415927)); // OK but does not make sense
or define new function double min(int, float);
```

C++ versus C linkage

- ◆ How can three different functions have the same name?
 - ◆ Look at what the compiler does


```
cd PT
cvs update -d
cd week3
g++ -c -save-temps -O3 min.C
```
 - ◆ Look at the assembly language file min.s and also at min.o


```
nm min.o
```
- ◆ The functions actually have different names!
 - ◆ Types of arguments appended to function name
- ◆ C and Fortran functions just use the function name
 - ◆ Can declare a function to have C-style name by using `extern "C"`

```
extern "C" { short min(short x, short y);}
```

Using macros (is dangerous)

- ◆ We still need many functions (albeit with the same name)
- ◆ In C we could use preprocessor macros:
 - ◆ `#define min(A,B) (A < B ? A : B)`
- ◆ However there are serious problems:
 - ◆ No type safety
 - ◆ Clumsy for longer functions
 - ◆ Unexpected side effects:


```
min(x++, y++); // will increment twice!!!
               // since this is: (x++ < y++ ? x++ : y++)
```
- ◆ Look at it:
 - ◆ `c++ -E minmacro.C`

Generic algorithms using templates in C++

- ◆ C++ templates allow a generic implementation:

```
template <class T>
inline T min (T x, T y)
{
    return (x < y ? x : y);
}
```

$$\min(x,y) \text{ is } \begin{cases} x & \text{if } x < y \\ y & \text{otherwise} \end{cases}$$

- ◆ Using templates we get functions that
 - ◆ work for many types T
 - ◆ are optimal and efficient since they can be inlined
 - ◆ are as generic and abstract as the formal definition
 - ◆ are one-to-one translations of the abstract algorithm

Usage Causes Instantiation

```
template <class T>
T min(T x, T y)
{
    return x < y ? x : y;
}
```

```
int x = min(3, 5);
int y = min(x, 100);
```

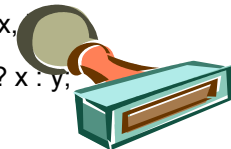
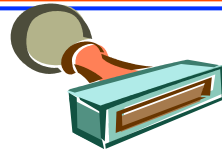
```
float z = min(3.14159f, 2.7182f);
```

// T is int

```
int min<int>(int x,
{
    return x < y ? x : y;
}
```

// T is float

```
float min<float>(float x, float y)
{
    return x < y ? x : y;
}
```



Discussion

“What is Polymorphism?”

Our definition:

Using many different types through the same interface

Generic programming process

- ◆ Identify useful and efficient algorithms
- ◆ Find their generic representation
 - ◆ Categorize functionality of some of these algorithms
 - ◆ What do they need to have in order to work **in principle**
- ◆ Derive a set of (minimal) requirements that allow these algorithms to run (efficiently)
 - ◆ Now categorize these algorithms and their requirements
 - ◆ Are there overlaps, similarities?
- ◆ Construct a framework based on classifications and requirements
- ◆ Now realize this as a software library

Generic Programming Process: Example

- ◆ (Simple) Family of Algorithms: min, max
- ◆ Generic Representation

$$\min(x,y) = \begin{cases} x & \text{if } x < y \\ y & \text{otherwise} \end{cases}$$

$$\max(x,y) = \begin{cases} x & \text{if } x > y \\ y & \text{otherwise} \end{cases}$$

- ◆ Minimal Requirements?
- ◆ Find Framework: Overlaps, Similarities?

Generic Programming Process: Example

- ◆ (Simple) Family of Algorithms: min, max
- ◆ Generic Representation

$$\min(x,y) = \begin{cases} x & \text{if } x < y \\ y & \text{otherwise} \end{cases}$$

$$\max(x,y) = \begin{cases} x & \text{if } y < x \\ y & \text{otherwise} \end{cases}$$

- ◆ Minimal Requirements yet?
- ◆ Find Framework: Overlaps, Similarities?

Generic Programming Process: Example

◆ Possible Implementation

```
template <class T>
T min(T x, T y)
{
    return x < y ? x : y;
}
```

◆ What are the Requirements on T?

- ◆ operator < , result convertible to bool

Generic Programming Process: Example

◆ Possible Implementation

```
template <class T>
T min(T x, T y)
{
    return x < y ? x : y;
}
```

◆ What are the Requirements on T?

- ◆ operator < , result convertible to bool
- ◆ Copy construction: need to copy the result!

Generic Programming Process: Example

◆ Possible Implementation

```
template <class T>
T const& min(T const& x, T const& y)
{
    return x < y ? x : y;
}
```

◆ What are the Requirements on T?

- ◆ operator < , result convertible to bool
- ◆ that's all!

The problem of different types: manual solution

◆ What if we want to call min(1,3.141)?

```
template <class R,U,T>
R const& min(U const& x, T const& y)
{
    return static_cast<R>(x < y ? x : y);
}
```

◆ Now we need to specify the first argument since it cannot be deduced.

```
min<double>(1,3.141);
min<int>(3,4);
```


Concepts

- ◆ A concept is a set of requirements, consisting of valid expressions, associated types, invariants, and complexity guarantees.
- ◆ A type that satisfies the requirements is said to model the concept.
- ◆ A concept can extend the requirements of another concept, which is called refinement.
- ◆ A concept is completely specified by the following:
 - ◆ Associated Types: The names of auxiliary types associated with the concept.
 - ◆ Valid Expressions: C++ expressions that must compile successfully.
 - ◆ Expression Semantics: Semantics of the valid expressions.
 - ◆ Complexity Guarantees: Specifies resource consumption (e.g., execution time, memory).
 - ◆ Invariants: Pre and post-conditions that must always be true.

Assignable concept

- ◆ Notation
 - ◆ X A type that is a model of Assignable
 - ◆ x, y Object of type X

Expression	Return type	Semantics	Postcondition
x=y;	X&	Assignment	X is equivalent to y
swap(x,y)	void	Equivalent to { X tmp = x; x = y; y = tmp; }	

CopyConstructible concept

◆ Notation

- ◆ X A type that is a model of CopyConstructible
- ◆ x, y Object of type X

Expression	Return type	Semantics	Postcondition
X(y)	X&		Return value is equivalent to y
X x(y);		Same as X x; x=y;	x is equivalent to y
X x=y;		Same as X x; x=y;	

Documenting a template function

◆ In addition to

- ◆ Preconditions
- ◆ Postconditions
- ◆ Semantics
- ◆ Exception guarantees

◆ The documentation of a template function must include

- ◆ Concept requirements on the types

◆ Note that the complete source code of the template function must be in a header file

Argument Dependent Lookup

- ◆ Also known as Koenig Lookup
- ◆ Applies only to *unqualified* calls
`abs(x)` ~~`std::abs(x)`~~
- ◆ Examines “associated classes and namespaces”
- ◆ Adds functions to overload set
- ◆ Originally for operators, e.g. `operator<<(std::ostream&, T);`



```

namespace lib {
    template <class T> T abs(T x)
    { return x > 0 ? x : -x; }

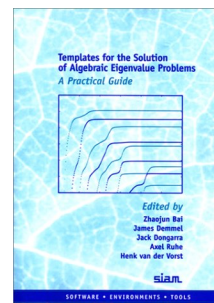
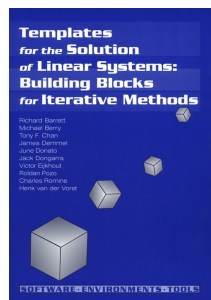
    template <class T>
    T compute(T x) {
        ...
        return abs(x);
    }
}

namespace user {
    class Num {};
    Num abs(Num);
    Num x = lib::compute(Num());
}
    
```

Diagram illustrating Argument Dependent Lookup (ADL) with arrows pointing from the `abs` call in `compute` and the `abs` call in `user` to the `abs` function in `lib` and the `abs` function in `user`.

Examples: iterative algorithms for linear systems

- ◆ Iterative template library (ITL)
 - ◆ Rick Lee *et al*, Indiana
- ◆ generic implementation of iterative solvers for linear systems from the “Templates” book
- ◆ Iterative Eigenvalue Template Library (IETL)
 - ◆ Prakash Dayal *et al*, ETH
- ◆ generic implementation of iterative eigensolvers. partially implements the eigenvalue templates book



The power method

- ◆ Is the simplest eigenvalue solver
 - ◆ returns the largest eigenvalue and corresponding eigenvector

```

ALGORITHM 4.1: Power Method for HEP
(1)   start with vector  $y = z$ , the initial guess
(2)   for  $k = 1, 2, \dots$ 
(3)      $v = y / \|y\|_2$ 
(4)      $y = Av$ 
(5)      $\theta = v^* y$ 
(6)     if  $\|y - \theta v\|_2 \leq \epsilon_M |\theta|$ , stop
(7)   end for
(8)   accept  $\lambda = \theta$  and  $x = v$ 
    
```

- ◆ Only requirements:
 - ◆ A is linear operator on a Hilbert space
 - ◆ Initial vector y is vector in the same Hilbert space
- ◆ Can we write the code with as few constraints?

Generic implementation of the power method

- ◆ A generic implementation is possible

```

OP A;
V v, y;
T theta, tolerance, residual;
...
do {
    v = y / two_norm(y);           // line (3)
    y = A * v;                     // line (4)
    theta = dot(v, y);             // line (5)
    v *= theta;                    // line (6)
    v -= y;
    residual = two_norm(v); // ||θ v - Av||
} while(residual > tolerance * abs(theta));
    
```

Concepts for the power method

- ◆ The triple of types (T,V,OP) models the Hilbertspace concept if
 - ◆ T must be the type of an element of a **field**
 - ◆ V must be the type of a vector in a **Hilbert space** over that field
 - ◆ OP must be the type of a **linear operator** in that Hilbert space

- ◆ All the allowed mathematical operations in a Hilbert space have to exist:
 - ◆ Let v, w be of type V
 - ◆ Let r, s of type T
 - ◆ Let a be of type OP.
 - ◆ The following must compile and have the same semantics as in the mathematical concept of a Hilbert space:
 - $r+s, r-s, r/s, r*s, -r$ have return type T
 - $v+w, v-w, v*r, r*v, v/r$ have return type V
 - $a*v$ has return type V
 - $two_norm(v)$ and $dot(v,w)$ have return type T
 - ...
 - ◆ Exercise: complete these requirement